

GPS



- Bc637PCIe (GPS)

One of the most used GPS functionality is the possibility of synchronization.

- ✓ Synchronization of the system clock with UTC (it is possible and quite easy to create NTP or PTP server)
- ✓ Use interrupts to execute a code periodically (see next slides)

Synchronization in MicroBoone

In MicroBoone experiment the data are acquired in frames. Frames are 1.6 msec long with data sampled at 2MHz. But the DAQ clock runs at a frequency of 16 MHz.

We need to know the GPS time corresponding to the frame number NOW.

The effect is that we have in memory a 3-column table of GPS time, DAQ Clock Time, Frame Number, as many rows as seconds that we want to keep in a circular buffer.

What is an interrupt?

Interrupt is "the forced suspension of a program in execution, in order to run a routine associated to the event that have generated the interrupt"

let's try to express it more clearly

Type of interrupt

There are three types of interrupts:

- External
- Hardware
- Software (Supervisor Call)

Only External interrupts are interesting for our purpose, then I will talk only about this type of interrupts.

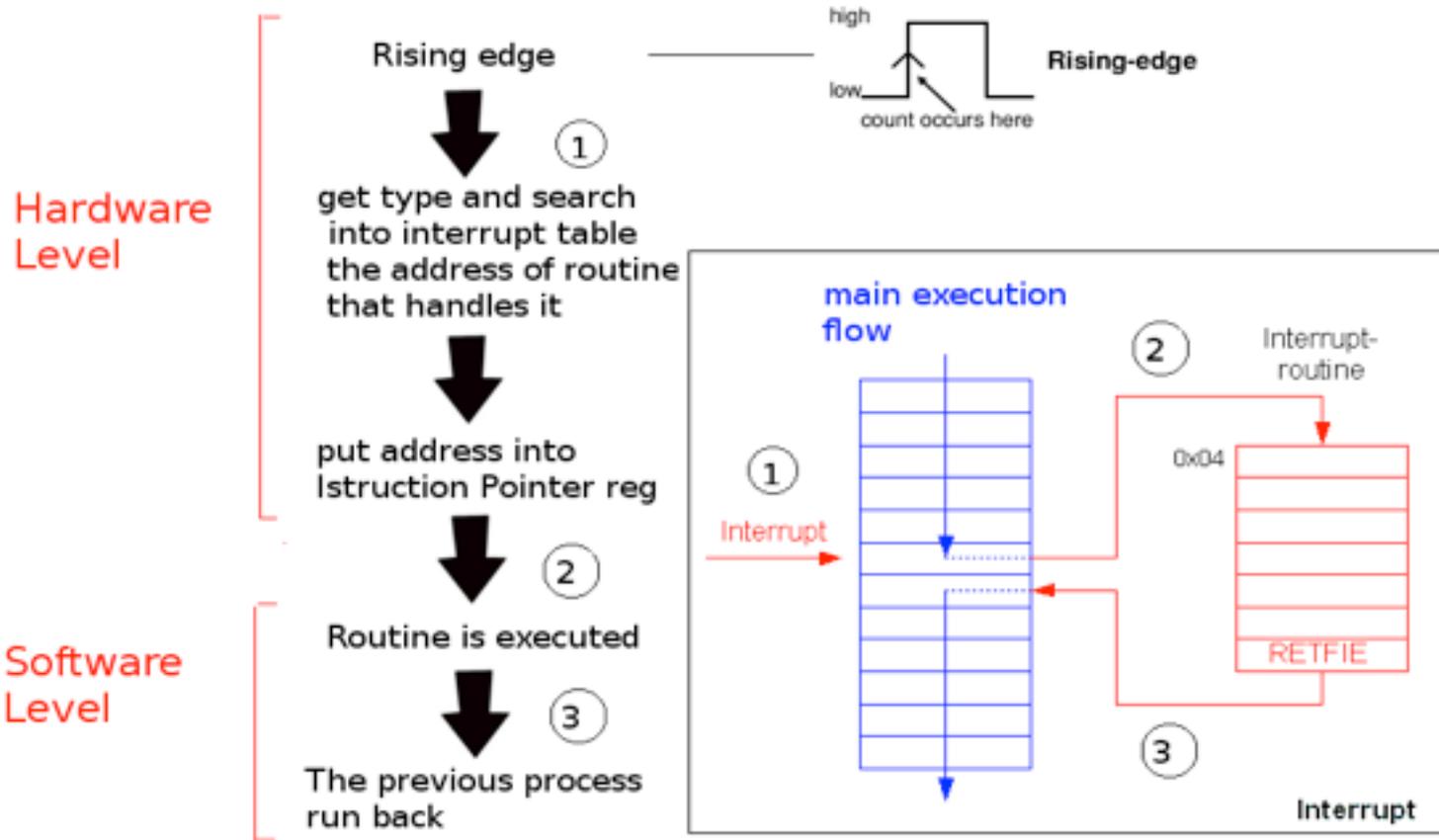
External Interrupts

This interrupts are associated with event that occurs outside CPUs.

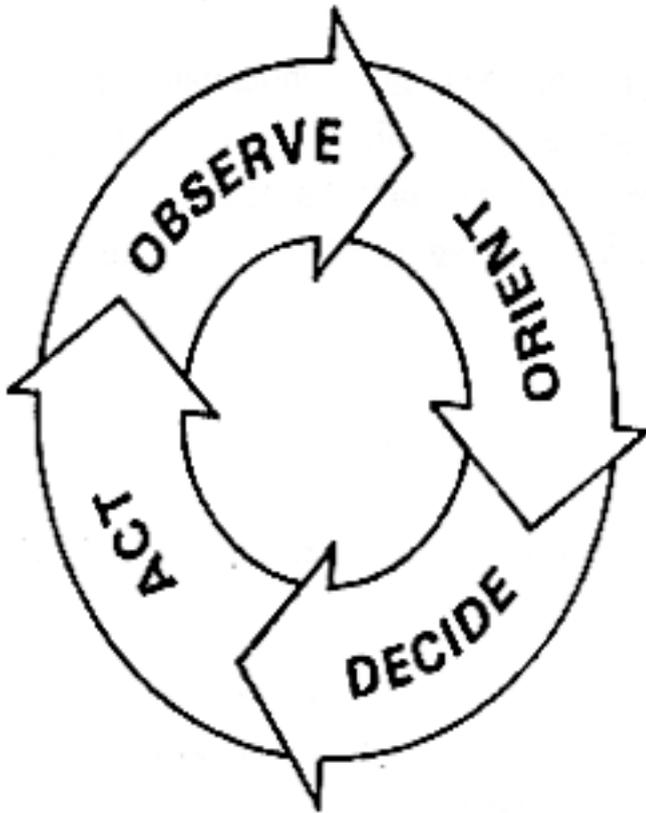
More concretely an events it is represented by a rising edge of an input signal.

In our application it corresponds with the pps signal that GPS card generates on PCI bus (it is a little bit more complex in reality)

Interrupt management flow



Why interrupts?



Interrupts are very useful for device management and this mechanism is the most used by modern operating systems (e.g. keyboard, mouse, screen are managed in this way).

For example they allow to avoid busy waiting.

Interrupt + GPS

- As we have already said GPS generate a periodic signal that is called Pulse Per Second
- It is possible to see this signal as an event, and to associate to it an appropriate routine.

Interrupt + GPS 2

- In this way we can obtain a periodic routine, with a period accurate at nanoseconds
- It is possible to implements periodic routine also using timer (I've done it in the consumer, see later)

Difference PPS/Timer

- There are two important advantages in use PPS+interrupts implementation
 - PPS is strongly more accurate than Timer
 - PPS is Synchronized with Coordinated Universal Time.
Therefore, for example, two different hosts, using PPS, execute their routine at the same time.

My codes

I've wrote some sample codes that implement this mechanism, called fancifully:

- GPS_1
- GPS_1_5
- GPS_2
- connection_control

GPS_1

This code include a main that sets up the card and the OS Interrupt table, and an interrupt routine, associated with PPS event, that reads directly time and print it.

The role of this is just demonstrate how to write and set a periodic PPS code.

GPS_1 code: main

Main calls initialization functions and wait:

- `bcStartPci()`; sets and starts the device
- `pci_set_ints(hBC_PCI)`; sets Interrupt Descriptor Table placing handler routine in it (`bcStartIntEx(hBC_PCI, bcIntHandlerRoutine, INTERRUPT_1PPS & 0x7F)`)

GPS_1 code: interrupt

Interrupt code:

```
//Get current time
```

```
//bcReadDecTimeEx is a library function that reads current time
```

```
bcReadDecTimeEx (hBC_PCI, &dectime, &min, &nano, &stat);
```

```
//Print time read
```

```
printf("pps_routine: %02d:%02d:%02d.%06lu%d \n",
```

```
    dectime.tm_hour, dectime.tm_min, dectime.tm_sec, min,  
    nano);
```

GPS_1 output

OUTPUT:

```
...  
bcIntHandlerRoutine  
pps_routine: 80:156:66.0040402  
bcIntHandlerRoutine  
pps_routine: 80:156:67.0140657  
bcIntHandlerRoutine  
pps_routine: 80:156:68.0030917  
bcIntHandlerRoutine  
pps_routine: 80:156:69.0041193  
bcIntHandlerRoutine  
pps_routine: 80:156:70.0041424  
bcIntHandlerRoutine  
pps_routine: 80:156:71.0041674  
bcIntHandlerRoutine  
pps_routine: 80:156:72.0141917  
bcIntHandlerRoutine  
pps_routine: 80:156:73.0052188  
bcIntHandlerRoutine  
pps_routine: 80:156:74.0032431  
bcIntHandlerRoutine  
pps_routine: 80:156:75.0032680  
bcIntHandlerRoutine  
pps_routine: 80:156:76.0022934  
...
```

This output demonstrate that:

- Routine occurs each second
- Time is captured with a random delay, due by the execution of software between interrupt signal and data acquisition.

GPS_1_5

This code is similar to the previous (GPS_1). The difference lies in how the current time is obtained: in this new implementation we use event register, latching (on hardware level) the time in which PPS arrives in it.

This code demonstrate how to latch an event (that could be different from PPS) time

Code differences between 1 and 1_5

Into main I need to setup event register:

```
iVal=1;  
EvDat.evtsrc = (BYTE)iVal;  
EvDat.evtctl = (BYTE)iVal;  
  
iVal=0;  
EvDat.evtlock = (BYTE)iVal;  
EvDat.evtsense = (BYTE)iVal;  
bcSetEventsData (hBC_PCI, &EvDat);
```

Into interrupt routine I only need to take time with
another function:

```
bcReadEventTimeEx (hBC_PCI, &evtmaj, &evtmin, &evtnano, &stat);
```

GPS_1_5 output

OUTPUT:

```
...  
Time: 09/13/2012 17:01:07.0000000 Status: 7  
Time: 09/13/2012 17:01:08.0000000 Status: 7  
Time: 09/13/2012 17:01:09.0000000 Status: 7  
Time: 09/13/2012 17:01:10.0000000 Status: 7  
Time: 09/13/2012 17:01:11.0000000 Status: 7  
Time: 09/13/2012 17:01:12.0000000 Status: 7  
Time: 09/13/2012 17:01:13.0000000 Status: 7  
Time: 09/13/2012 17:01:14.0000000 Status: 7  
Time: 09/13/2012 17:01:15.0000000 Status: 7  
Time: 09/13/2012 17:01:16.0000000 Status: 7  
Time: 09/13/2012 17:01:17.0000000 Status: 7  
...
```

N.B.: This output does not demonstrate anything about precision of GPS: This is GPS time, and not UTC time. Then is obvious that GPS says that signal, that it thinks to generate each second, has a infinite precision.

GPS_2

This code evolves GPS_1 (it doesn't use event register), implementing the classic paradigm of communication Producer-Consumer, where producer is the interrupt and consumer is a stand-alone thread. The shared information is the number of PPS that are arrived.

Shared memory

- In computing, shared memory is memory that may be simultaneously accessed by multiple programs with an intent to provide communication among them.
- Linux permits a controlled use of shared memory, showing it like a device called shm.

GPS_2 code

Producer(interrupt):

```
//each time I need to get shared
//address
memory_loc_key = ftok(".", 'M');
if((id = shmget( memory_loc_key,
sizeof(long int), 0666 ))<0){/*...*/ }
shm_ptr = shmat (id,NULL,0);
if((int)shm_ptr == -1){...}

//increment pps_counter
(*(long int*)shm_ptr)++;
```

Consumer(thread):

```
//get shared address (one time)
memory_loc_key = ftok(".", 'M');
if((id = shmget( memory_loc_key, sizeof
(long int), 0666 ))<0){/*...*/ }
shm_ptr = shmat (id,NULL,0);
if((int)shm_ptr == -1){...}
[...]
While(1){
    //use data
    printf("I'm the consumer!: %ld \n", *((long
int*)shm_ptr));
    /*wait 0.5 sec*/
}
```

GPS_2 output

```
...  
bcIntHandlerRoutine  
I'm the consumer!: 2926  
I'm the consumer!: 2926  
pps_routine: 87:144:244.0059356  
bcIntHandlerRoutine  
I'm the consumer!: 2927  
I'm the consumer!: 2927  
pps_routine: 87:144:245.0039618  
bcIntHandlerRoutine  
I'm the consumer!: 2928  
I'm the consumer!: 2928  
pps_routine: 87:144:246.0059860  
bcIntHandlerRoutine  
I'm the consumer!: 2929  
I'm the consumer!: 2929  
pps_routine: 87:144:247.0030130  
...
```

connection_control

This simple code periodically (about 4 sec) requires packet46 to GPS card, and reads in it information about status of connection. Then it prints a message describing this status.

connection_control output

...
Status: No usable satellites
...

This output was taken
when the antenna
was disconnected.

Conclusions

The codes that I've briefly explained are simply sample codes, and they try to show functionalities. It would be simple to modify them, in order to satisfy the various needs.